



Presshammer

Rowhammer and Rowpress without Physical Address Information

Jonas Juffinger **Sudheendra Raghav Neela** **Martin Heckel** **Lukas Schwarz**
Florian Adamsky **Daniel Gruss**

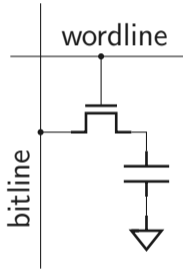
2024-07-19

IAIK, Graz University of Technology, Austria & Hof University of Applied Sciences, Germany

DIMVA 2024 – EPFL, Lausanne, Switzerland

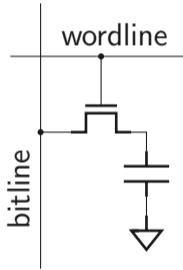
Introduction

Dynamic Random Access Memory (DRAM)

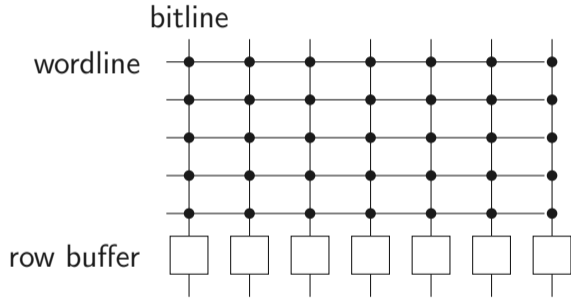


DRAM Cell

Dynamic Random Access Memory (DRAM)

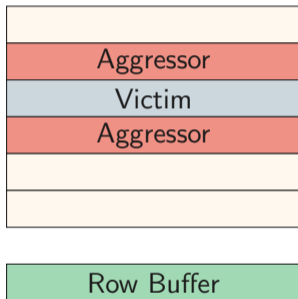


DRAM Cell



DRAM Bank

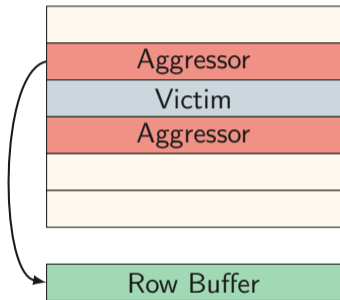
Rowhammer [2]



Repeatedly and **quickly** active multiple rows.

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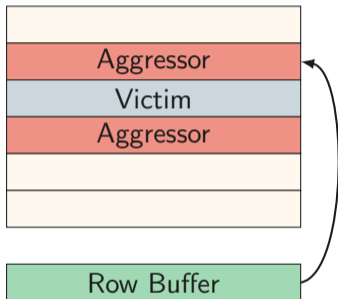
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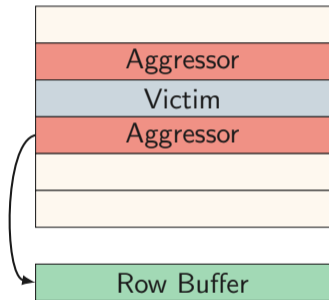
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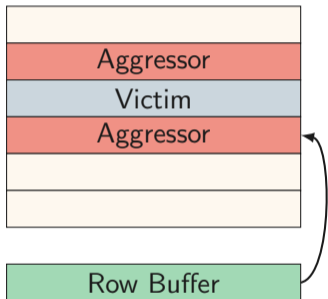
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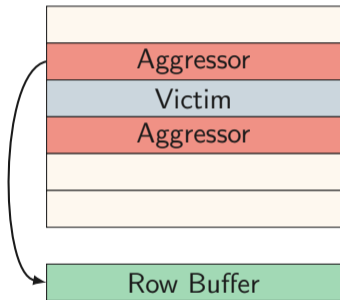

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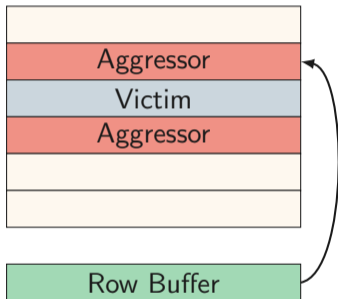
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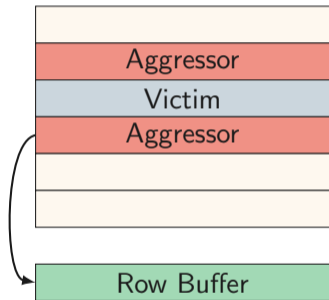
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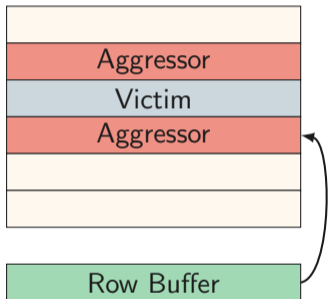
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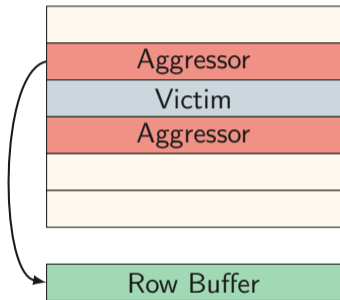
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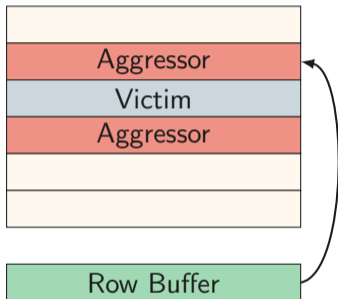
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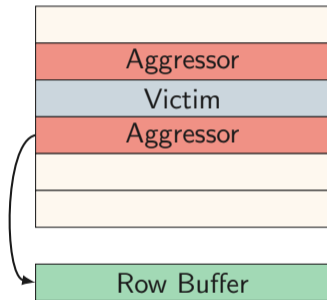
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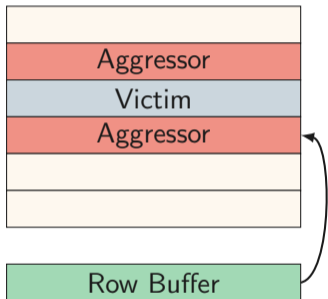
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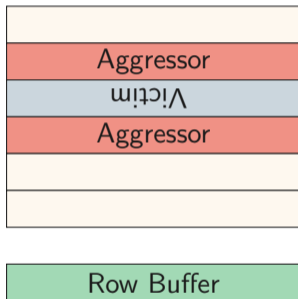

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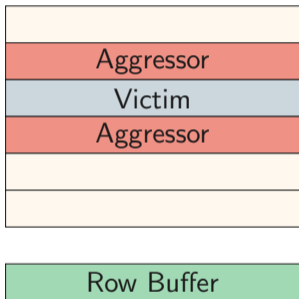
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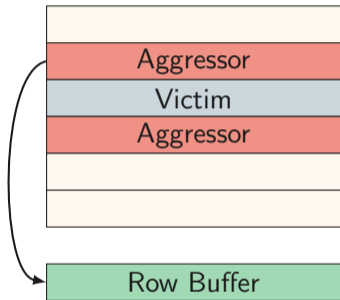
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Keep rows open as **long** as possible.



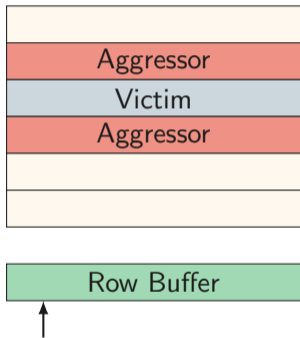
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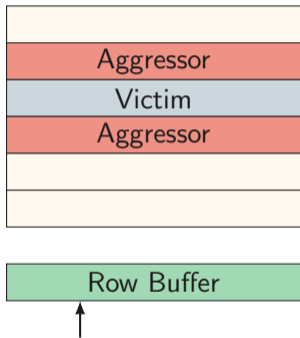
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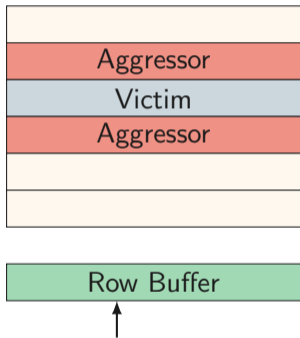
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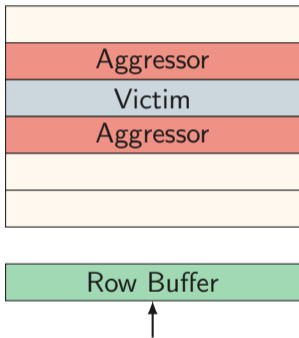
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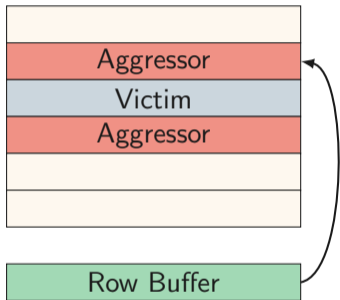
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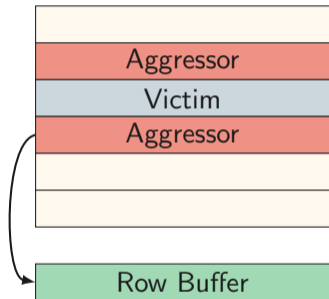
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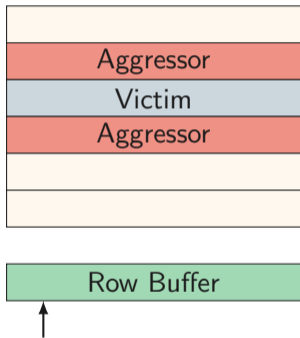
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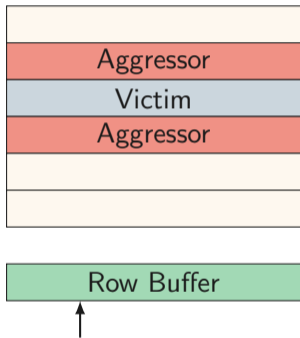
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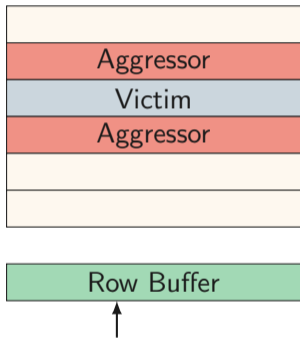
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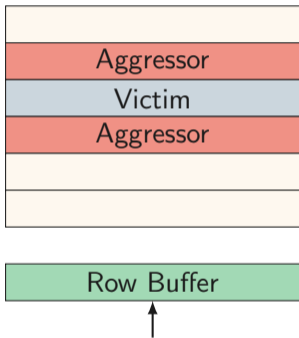
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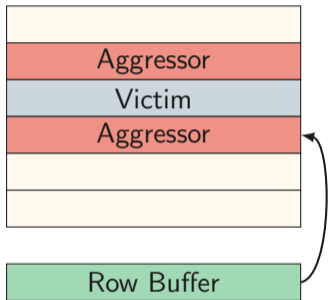
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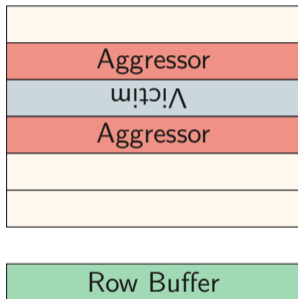


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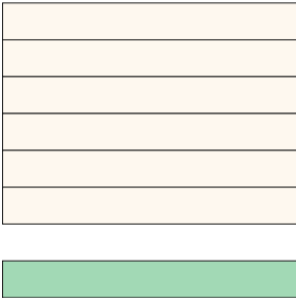


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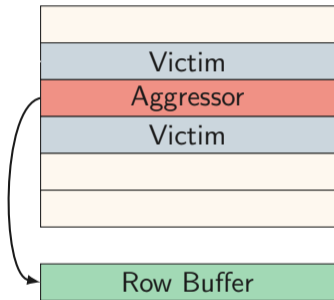
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Flips different bits than Rowhammer

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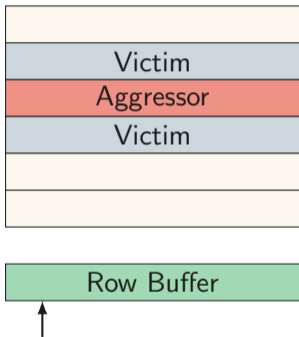
Keep **one** row open as long as possible.

Single-Sided RowPress



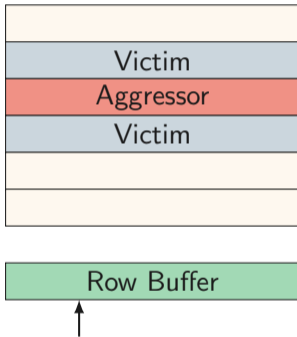
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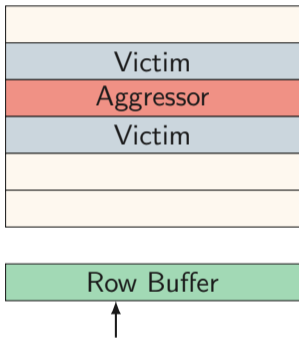
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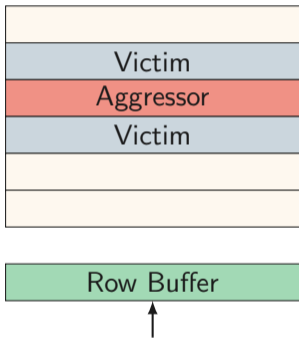
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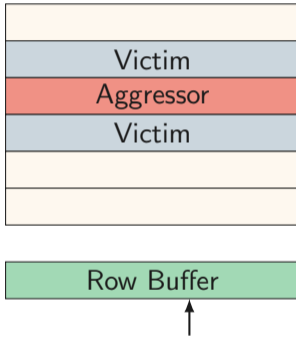
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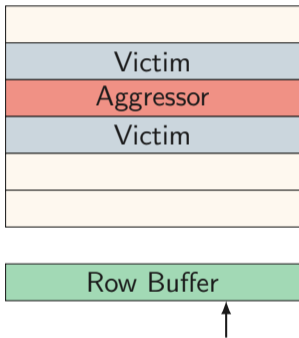
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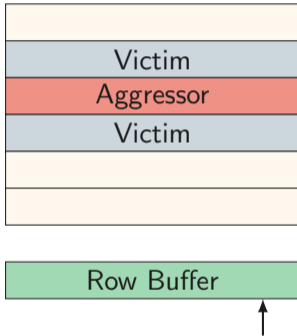
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Single-Sided RowPress



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Single-Sided RowPress



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But wait...



There is this paper from 2018.

Another Flip in the Wall of Rowhammer Defenses

Daniel Gruss¹, Moritz Lipp¹, Michael Schwarz¹, Daniel Genkin²,
Jonas Juffinger¹, Sioli O’Connell³, Wolfgang Schoechl¹, and Yuval Yarom^{3,4}

¹ Graz University of Technology

² University of Pennsylvania and University of Maryland

³ University of Adelaide

⁴ Data61

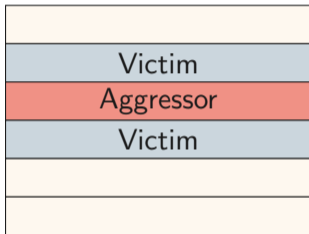
Abstract—The Rowhammer bug allows unauthorized modification of bits in DRAM cells from unprivileged software, enabling powerful privilege-escalation attacks. Sophisticated Rowhammer countermeasures have been presented, aiming at mitigating the Rowhammer bug or its exploitation. However, the state of the art provides insufficient insight on the completeness of these defenses.

In this paper, we present novel Rowhammer attack and exploitation primitives, showing that even a combination of all defenses is ineffective. Our new attack technique, *one-location hammering*, breaks previous assumptions on requirements for triggering the Rowhammer bug, i.e., we do not hammer multiple DRAM rows but only keep one DRAM row constantly open. Our new exploitation technique, *opcode flipping*, bypasses recent isolation mechanisms by flipping bits in a predictable and targeted way in userspace binaries. We replace conspicuous

Software-based mitigations, which can be implemented on commodity systems, have also been proposed. These include ad-hoc defense techniques such as doubling the RAM refresh rates [44], removing unprivileged access to the pagemap interface [45, 62, 65], and prohibiting the `clflush` instruction [65]. However, recent works have already bypassed these countermeasures [6, 24, 68]. Other ad-hoc attempts, such as disabling page deduplication by default [52, 60], only prevent specific Rowhammer attacks exploiting these features [11, 59], but not all Rowhammer attacks.

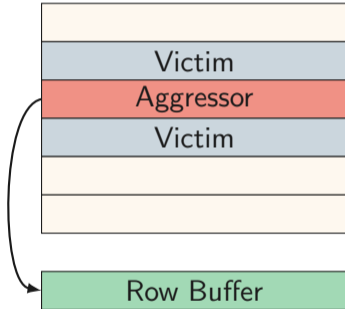
The research community proposed sophisticated defenses which seemingly have solved the Rowhammer problem. Based on the underlying primitives of these defenses we introduce

One-Location Rowhammer [1]



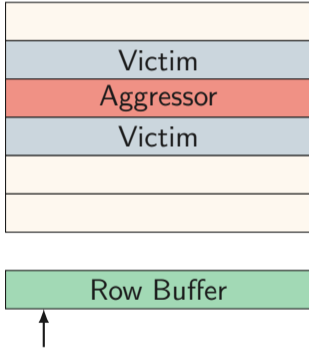
Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



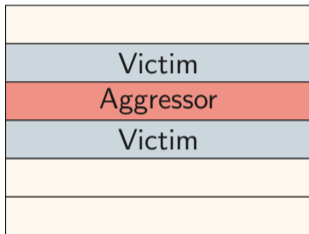
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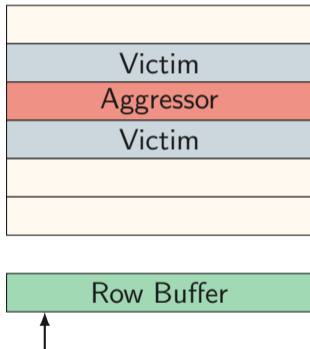
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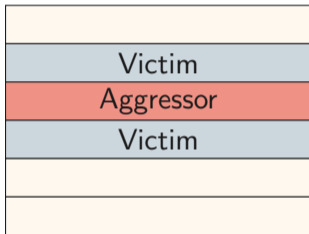
Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



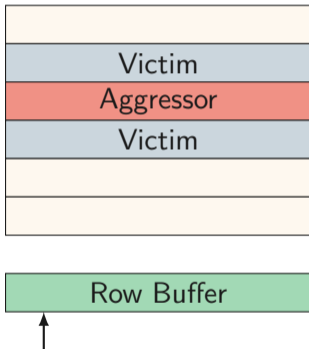
Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



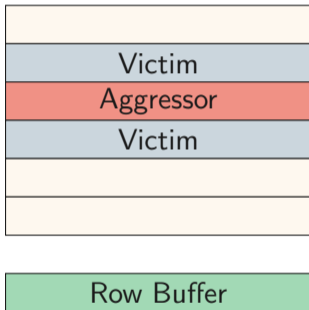
Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



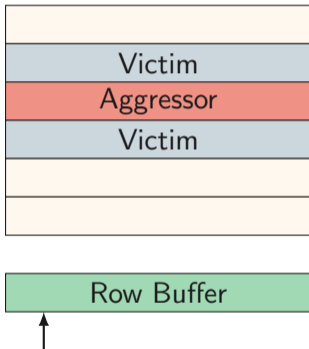
Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



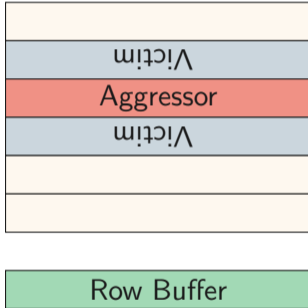
Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



Repeatedly and **quickly** activate **one** row.

One-Location Rowhammer [1]



Repeatedly and **quickly** activate **one** row.



LOOKS A LOT LIKE ROWPRESS

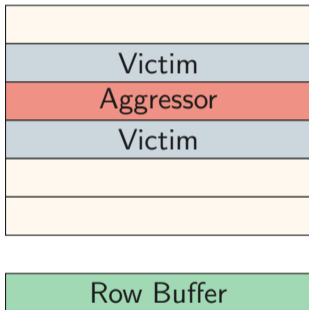
One-Location Rowhammer - Hypothesis





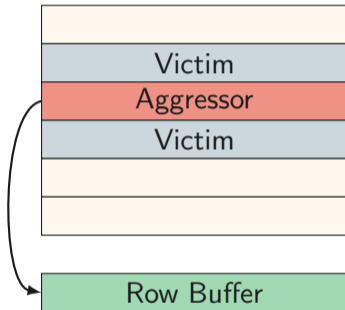
It's the memory controller.

One-Location Rowhammer - Hypothesis



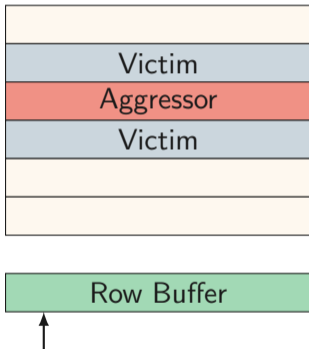
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



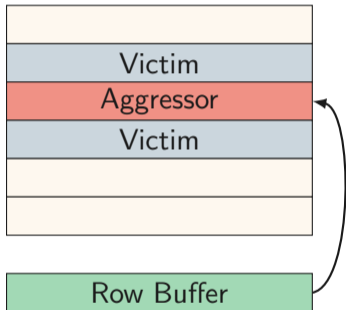
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



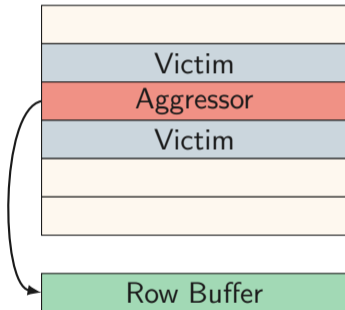
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



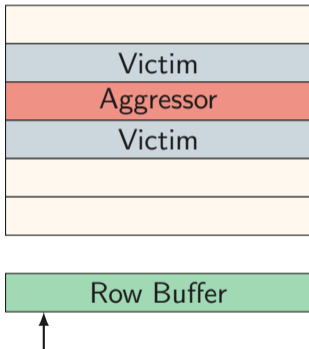
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



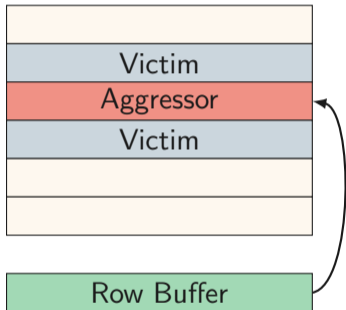
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



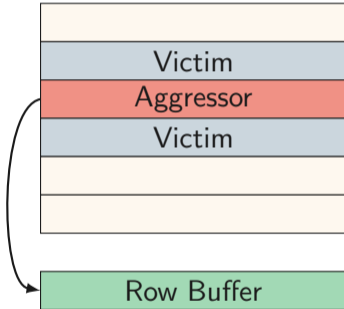
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



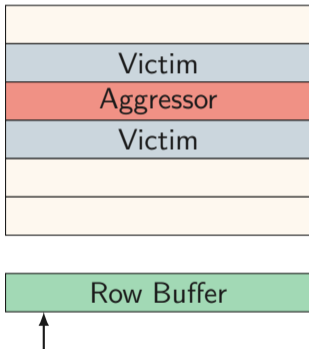
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



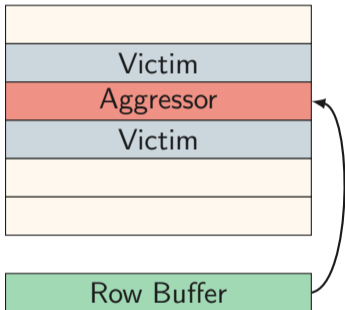
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



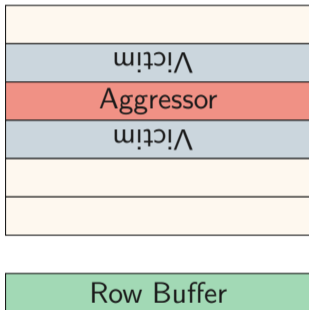
Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



Memory controller employs a *closed-row* policy.

One-Location Rowhammer - Hypothesis



Memory controller employs a *closed-row* policy.

Case closed ✓

Is it?

Is Rowpress = OL Rowhammer?

Is Rowpress = OL Rowhammer?

or where Gruss et al. wrong all along?



RH and RP are caused by different physical effects

RH and RP are caused by different physical effects



Different weak cells i.e. bit flip locations

RH and RP are caused by different physical effects

Different weak cells i.e. bit flip locations

Bit flip locations of OL RH = Rowpress?

RH and RP are caused by different physical effects

Different weak cells i.e. bit flip locations

Bit flip locations of OL RH = Rowpress?

YES

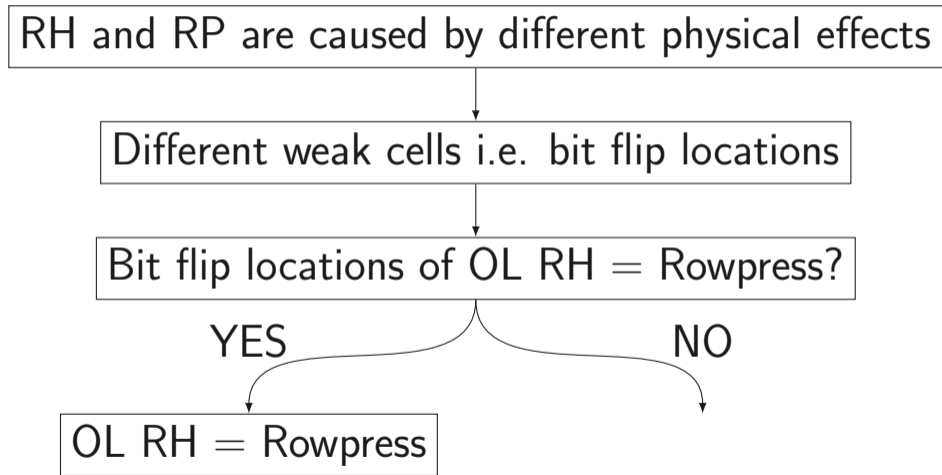
RH and RP are caused by different physical effects

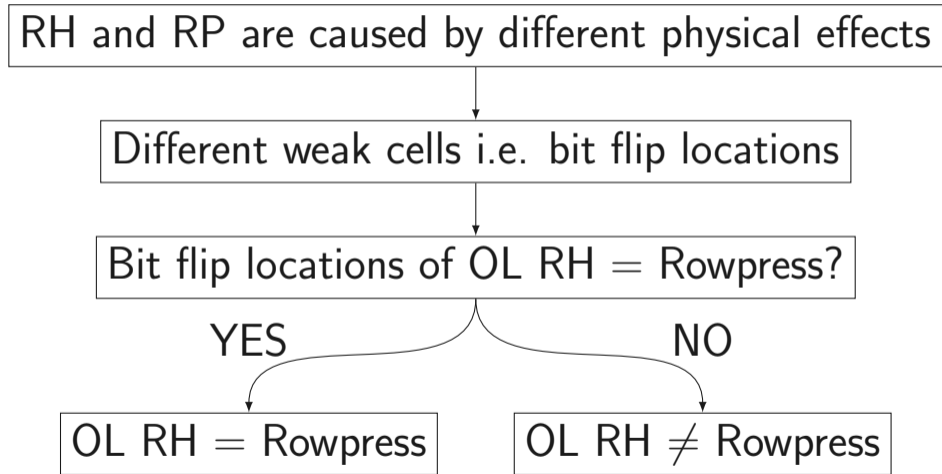
Different weak cells i.e. bit flip locations

Bit flip locations of OL RH = Rowpress?

YES

OL RH = Rowpress





		RH		RP	
		DS	OL	DS	OL
RH	DS				
	OL				
RP	DS				
	OL				

	DIMM	Size	MT/s	Double Sided		One Location	
				RH*	RP	RH	RP
DDR4	D1	8 GB	2133	✓	✓	✓	✓
	D2	8 GB	2133	✓	✓	✓	✓
	D3	8 GB	2400	✓	✗	~	~
	D4	8 GB	2400	✓	✗	~	~
	D5	8 GB	2133	✗	✗	~	~
	D6	8 GB	2666	✓	✗	~	~
	D7	8 GB	2666	✗	✗	~	~
	D8	8 GB	2666	✗	✗	~	~
	D9	8 GB	2666	✓	✗	~	~
	D10	8 GB	2666	✗	✗	~	~
	D11	8 GB	3000	✗	✗	~	~
	D12	8 GB	2133	✗	✗	~	~

		RH		RP	
		DS	OL	DS	OL
RH	DS	100	56.5	13.7	1.0
	OL	61.8	100	14.5	1.1
RP	DS	98.1	95.3	100	1.0
	OL	96.3	98.8	81.4	100

PressHammer Exploit



Exploit Steps

- Find complete DRAM row



Exploit Steps

- Find complete DRAM row
- Template memory for bit flips



Exploit Steps

- Find complete DRAM row
- Template memory for bit flips
- Spray page tables



Exploit Steps

- Find complete DRAM row
- Template memory for bit flips
- Spray page tables
- Flip page tables



Exploit Steps

- Find complete DRAM row
- Template memory for bit flips
- Spray page tables
- Flip page tables
- → Profit

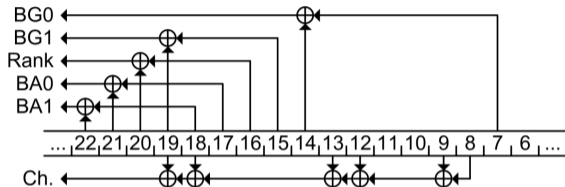
- DRAM row typically 4 kB

Find Complete DRAM Row

- DRAM row typically 4 kB
- Memory Page typically 4 kB

Find Complete DRAM Row

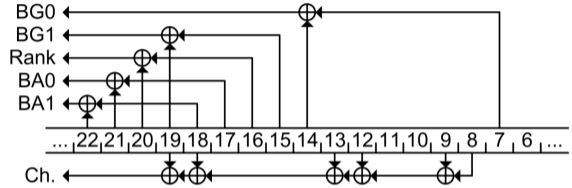
- DRAM row typically 4 kB
- Memory Page typically 4 kB
- No 1:1 mapping



Pessl et al. [4]

Find Complete DRAM Row

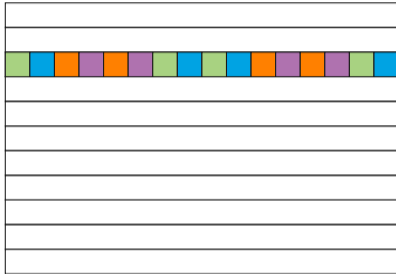
- DRAM row typically 4 kB
- Memory Page typically 4 kB
- No 1:1 mapping



Pessl et al. [4]



DRAM Row Finding



Bank 0

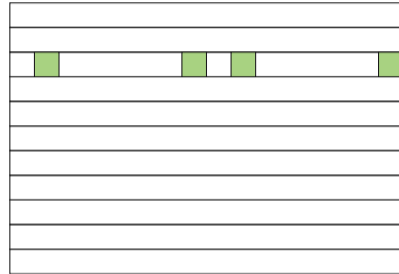


Bank 1

DRAM Row Finding



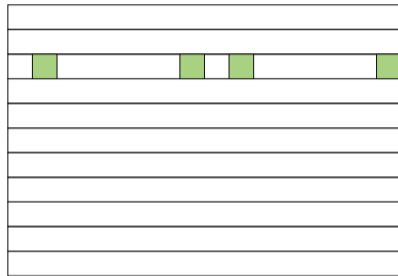
Bank 0



Bank 1

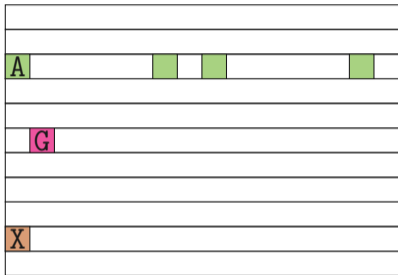


Bank 0

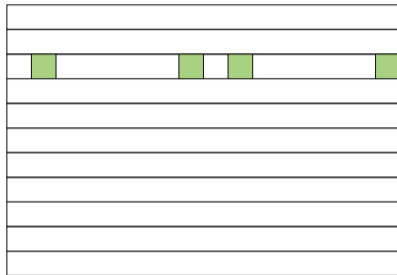


Bank 1

R1 Bank conflict with X



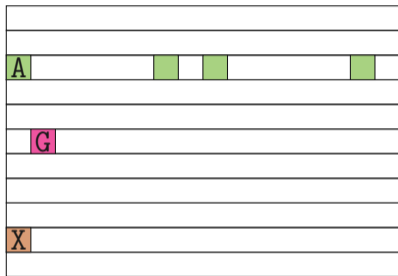
Bank 0



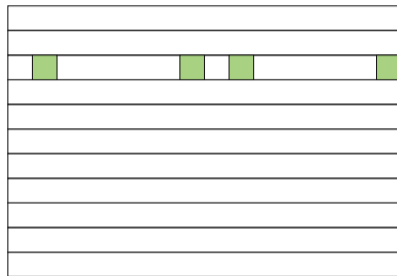
Bank 1

R1 Bank conflict with X

DRAM Row Finding



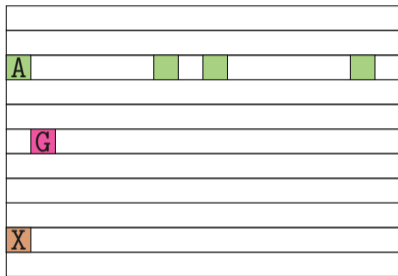
Bank 0



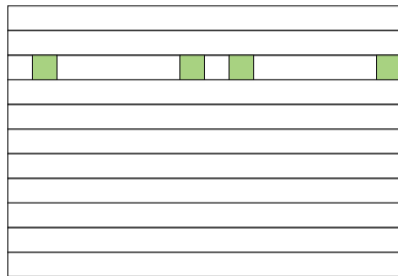
Bank 1

R1 Bank conflict with X ✓

DRAM Row Finding



Bank 0

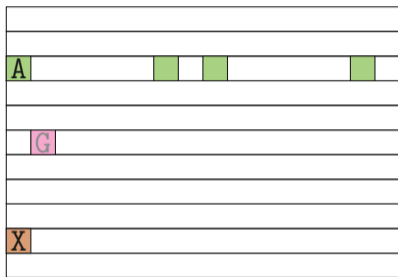


Bank 1

R1 Bank conflict with X ✓

R2 No bank conflict with A

DRAM Row Finding



Bank 0

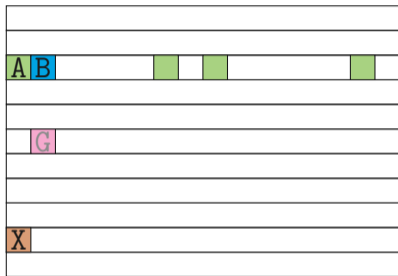


Bank 1

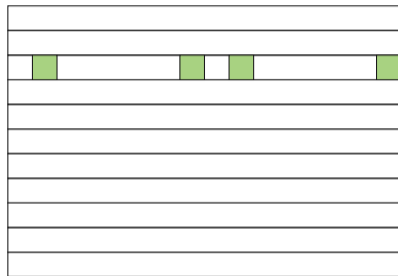
R1 Bank conflict with X ✓

R2 No bank conflict with A ✗

DRAM Row Finding



Bank 0

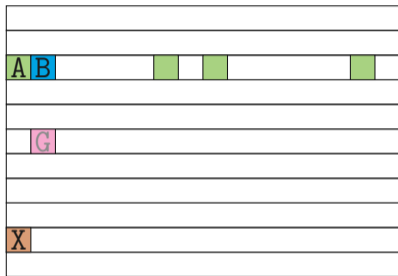


Bank 1

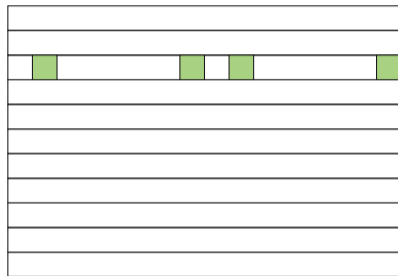
R1 Bank conflict with X

R2 No bank conflict with A

DRAM Row Finding



Bank 0



Bank 1

R1 Bank conflict with X ✓

R2 No bank conflict with A

DRAM Row Finding



Bank 0

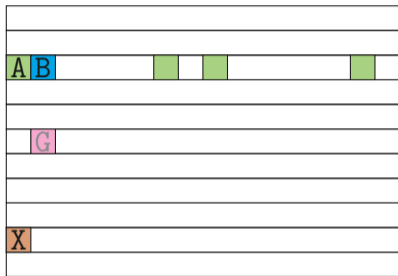


Bank 1

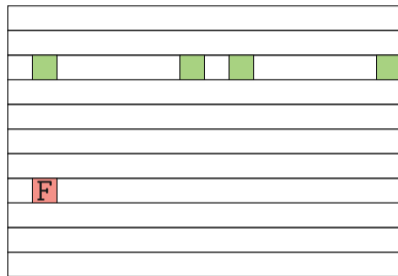
R1 Bank conflict with X ✓

R2 No bank conflict with A ✓

DRAM Row Finding



Bank 0

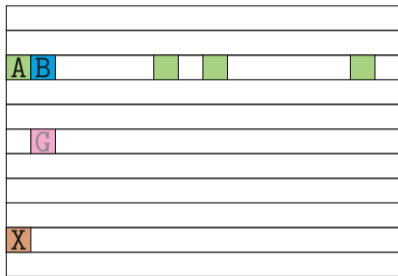


Bank 1

R1 Bank conflict with X

R2 No bank conflict with A

DRAM Row Finding



Bank 0

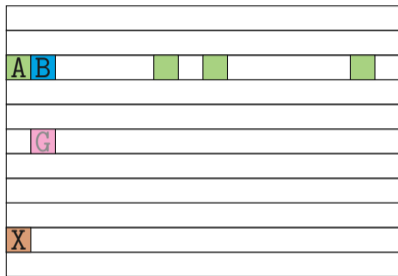


Bank 1

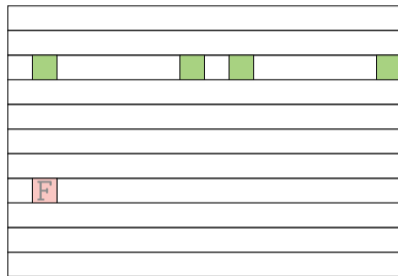
R1 Bank conflict with X ✗

R2 No bank conflict with A

DRAM Row Finding



Bank 0

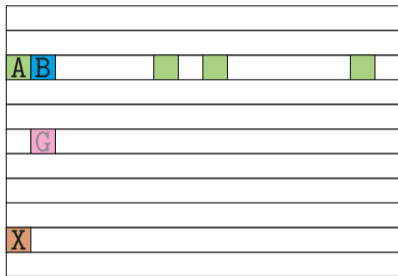


Bank 1

R1 Bank conflict with X ✗

R2 No bank conflict with A

DRAM Row Finding



Bank 0

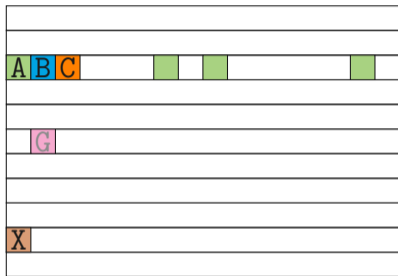


Bank 1

R1 Bank conflict with X ✗

R2 No bank conflict with A ✓

DRAM Row Finding



Bank 0

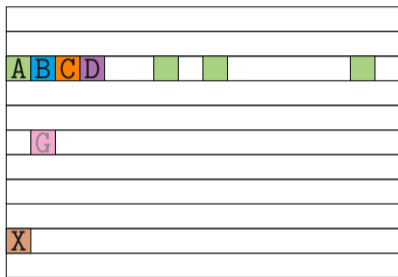


Bank 1

R1 Bank conflict with X ✓

R2 No bank conflict with A ✓

DRAM Row Finding



Bank 0



Bank 1

R1 Bank conflict with X ✓

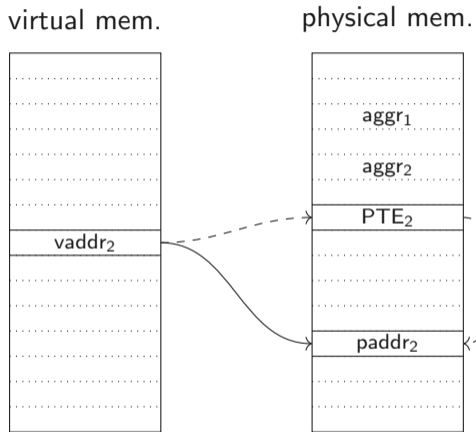
R2 No bank conflict with A ✓

DRAM Row Finding Accuracy and Speed

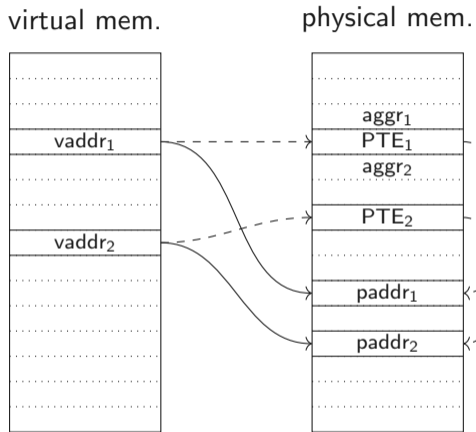


# Meas.	Duration per Row	Found CLs	Correct CLs
250	40 ms, $\sigma = 17$	124, $\sigma = 4.9$	119, $\sigma = 14$
500	45 ms, $\sigma = 15$	128, $\sigma = 0.8$	125, $\sigma = 16$
1000	116 ms, $\sigma = 40$	125, $\sigma = 4.2$	123, $\sigma = 4.3$
2500	260 ms, $\sigma = 109$	125, $\sigma = 4.9$	124, $\sigma = 5.2$

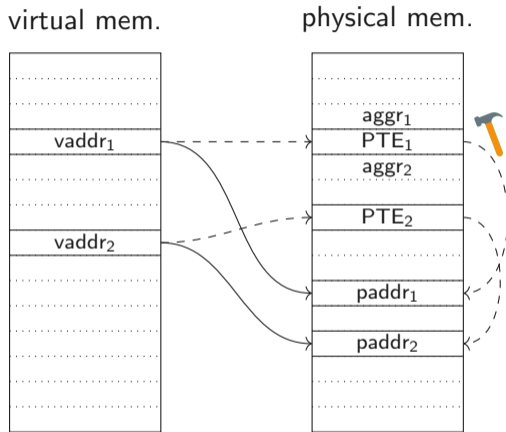
Page Table Entry Flipping



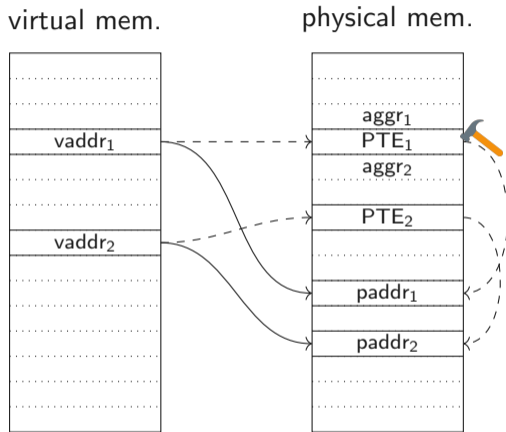
Page Table Entry Flipping



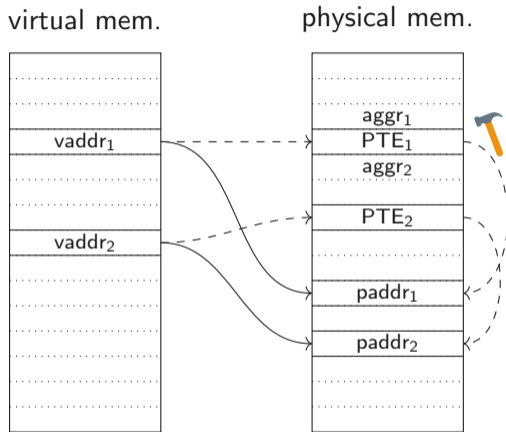
Page Table Entry Flipping



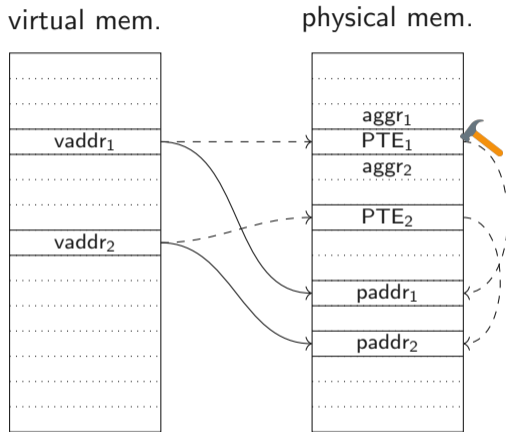
Page Table Entry Flipping



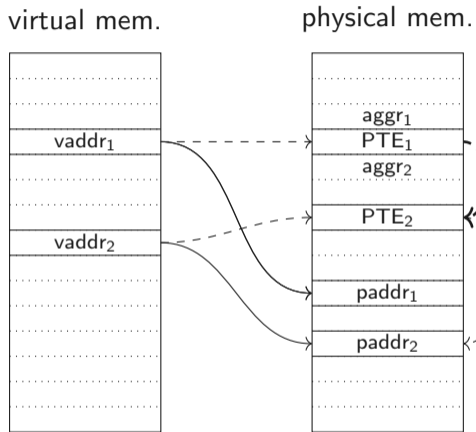
Page Table Entry Flipping



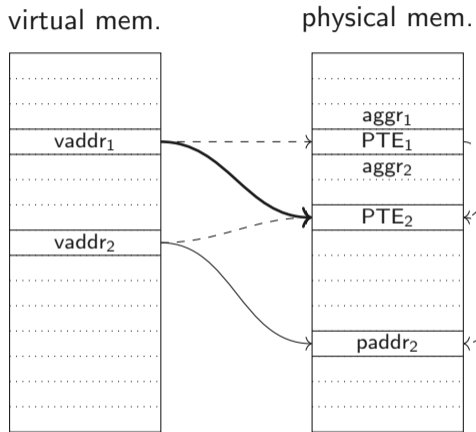
Page Table Entry Flipping



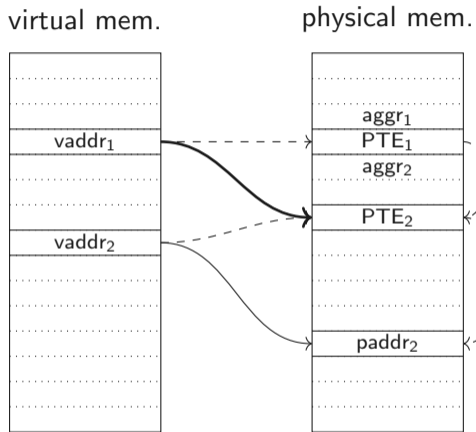
Page Table Entry Flipping



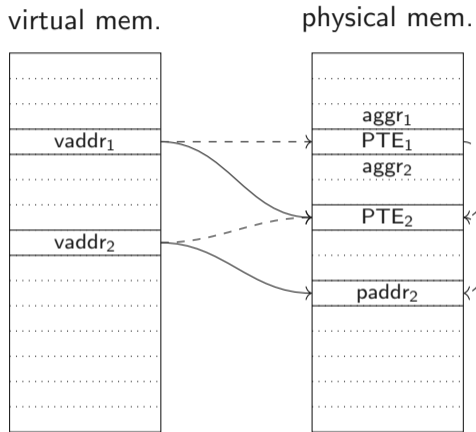
Page Table Entry Flipping



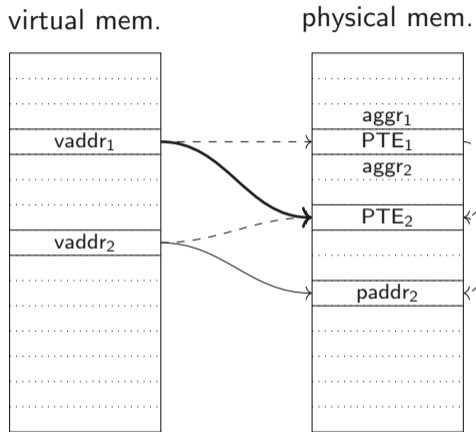
Page Table Entry Flipping



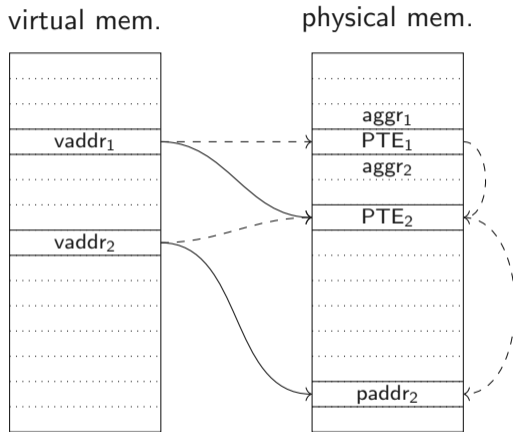
Page Table Entry Flipping



Page Table Entry Flipping



Page Table Entry Flipping







Did not successfully hammer page table.

DNF



585 s

Did not successfully hammer page table.

DNF



585 s



DNF

Did not successfully hammer page table.

Hammered page table, but didn't point to other.

DNF

Did not successfully hammer page table.

585 s

DNF

Hammered page table, but didn't point to other.

616 s

DNF

Did not successfully hammer page table.

585 s

DNF

Hammered page table, but didn't point to other.

616 s

337 s



Presshammer

Rowhammer and Rowpress without Physical Address Information

Jonas Juffinger Sudheendra Raghav Neela

Martin Heckel Lukas Schwarz

Florian Adamsky Daniel Gruss

2024-07-19

✉ jonas.juffinger@iaik.tugraz.at

🐦 [NotImaginary_](#)

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